

NOTE

ENUMERATION OF WORDS BY THEIR NUMBER OF MISTAKES

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Consider all words in $\{1, \dots, n\}$. A fixed set of words is labeled as the set of “mistakes”. A generating function for the number of words with m_1 1’s, \dots , m_n n ’s and k mistakes is given. This generalizes a result of Gessel who considered the case where all the mistakes are two-lettered. A similar result has been independently obtained by Goulden and Jackson.

1.

Fix an alphabet $\{1, \dots, n\}$. To every word $w = \sigma_1 \cdots \sigma_l$ we associate the monomial $x^w = x_{\sigma_1} \cdots x_{\sigma_l}$ in the non-commuting indeterminates x_1, \dots, x_n . A subword of $\sigma_1 \cdots \sigma_l$ is anything of the form $\sigma_i \sigma_{i+1} \cdots \sigma_j$, $1 \leq i \leq j \leq l$. Let L be a set of words to be labeled as “mistakes”. We assume that no proper subword of a mistake is a mistake. The number of subwords of w which belong to L is the number of mistakes of w and will be denoted by $d(w)$. For example if $L = \{123, 231\}$, $d(1231) = 2$, because both 123 and 231 belong to L . A word w is said to be of type (m_1, \dots, m_n) if it has m_1 1’s, m_2 2’s, \dots , m_n n ’s; e.g. the type of 12112331 is $(4, 2, 2)$. Let M be the set of words w such that every letter of w belongs to some mistake and every mistake, except the last, overlaps, on the right, with another mistake. For example if $L = \{123, 231, 312\}$, $M = \{123, 231, 312, 1231, 2312, 3123, 12312, 23123, 31231, \dots, \text{etc.}\}$.

The following is a generalization of Theorem 7.2 in Gessel [2]; Gessel’s theorem considers the case where L only contains two-lettered words.

Theorem

$$\sum_{w \in \text{all words}} t^{d(w)} x^w = \left[1 - x_1 - \cdots - x_n - \sum_{v \in M} (t-1)^{d(v)} x^v \right]^{-1}. \tag{1}$$

Proof. Let $s(w)$ denote the type of a word w . Let $C(\mathbf{m}) = C(m_1, \dots, m_n)$ be the set of words of type $\mathbf{m} = (m_1, \dots, m_n)$. Define

$$F(\mathbf{m}) = \sum_{w \in (\mathbf{m})} t^{d(w)} x^w. \tag{2}$$

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We shall prove that for $\mathbf{m} \neq \mathbf{0}$

$$F(\mathbf{m}) = \sum_{i=1}^n F(\mathbf{m} - \mathbf{e}_i)x_i + \sum_{v \in M} (t-1)^{d(v)}F(\mathbf{m} - s(v))x^v, \tag{3}$$

where $\mathbf{e}_i = (0, \dots, 1, 0, \dots, 0)$ with the 1 on the i th place.

This will be accomplished by showing that for any $w \in C(\mathbf{m})$, the coefficient of x^w in the r.h.s. of (3) is $t^{d(w)}$. Indeed, let w_2 be the maximal tail of w which belongs to M ; then $w = w_1w_2$ for some word w_1 , and $d(w) = d(w_1) + d(w_2)$. Note that w has $d(w_2)$ tails which belong to M and thus x^w appear $d(w_2) + 1$ times in the r.h.s of (3). Since w loses a mistake by chopping off its last letter and loses $k + 1$ mistakes by chopping off a tail which belongs to M and which has k mistakes, the coefficient of x^w in the r.h.s. of (3) is (Put $d(w_1) = d_1, d(w_2) = d_2$):

$$t^{d_1}[t^{d_2-1} + (t-1)t^{d_2-2} + (t-1)^2t^{d_2-3} + \dots + (t-1)^{d_2-1} + (t-1)^{d_2}] = t^{d_1}t^{d_2} = t^{d(w)}.$$

Here we used (2) with \mathbf{m} replaced by $\mathbf{m} - \mathbf{e}_i$ and $\mathbf{m} - s(v)$.

Let $\delta(\mathbf{m})$ be the discrete delta function: $\delta(\mathbf{0}) = 1; \delta(\mathbf{m}) = 0, \mathbf{m} \neq \mathbf{0}$. Then, since $F(\mathbf{0}) = 1$ and by convention F is zero outside \mathbf{N}^n :

$$F(\mathbf{m}) - \sum_{i=1}^n F(\mathbf{m} - \mathbf{e}_i)x_i - \sum_{v \in M} (t-1)^{d(v)}F(\mathbf{m} - s(v))x^v = \delta(\mathbf{m}).$$

Summing both sides over all $\mathbf{m} \in \mathbf{Z}^n$ yields

$$\left[\sum_{w \in \text{all words}} t^{d(w)}x^w \right] \left[1 - x_1 - x_2 - \dots - x_n - \sum_{v \in M} (t-1)^{d(v)}x^v \right] = 1,$$

from which (1) follows.

2. The commutative case

If we let x_1, \dots, x_n commute in (1) we obtain a generating function for $G(\mathbf{m}; k)$, the number of words of type \mathbf{m} with exactly k mistakes:

$$\sum G(m_1, \dots, m_n; k)x_1^{m_1} \dots x_n^{m_n} t^k = \left[1 - x_1 - \dots - x_n - \sum_{v \in M} (t-1)^{d(v)}x^v \right]^{-1}. \tag{4}$$

Example. $n = 3, L = \{123, 132\}$. Here $L = M$ and

$$\sum G(m_1, m_2, m_3, k)x_1^{m_1}x_2^{m_2}x_3^{m_3}t^k = [1 - x_1 - x_2 - x_3 + 2(1-t)x_1x_2x_3]^{-1}.$$

Putting $t = -1$ we get

$$\begin{aligned} &\text{coefficient of } x_1^{m_1}x_2^{m_2}x_3^{m_3} \text{ in } [1 - x_1 - x_2 - x_3 + 4x_1x_2x_3]^{-1} = \\ &\quad \#\{\text{words in } C(\mathbf{m}) \text{ with an even number of mistakes}\} \\ &\quad - \#\{\text{words in } C(\mathbf{m}) \text{ with an odd number of mistakes}\}. \end{aligned}$$

Askey and Gasper [1] proved that the l.h.s. is positive. It will be nice to give a direct proof that the r.h.s. is positive.

Finally let us mention that whenever L is finite but M is infinite it is still possible to evaluate the sum on the r.h.s. of (4) using the geometric series expansion of a certain matrix: $\sum A^k = (I - A)^{-1}$. Thus whenever L is finite the generating function $\sum G(\mathbf{m}; k) \mathbf{x}^{\mathbf{m}} t^k$ is a rational function. The details are left to the sufficiently interested reader.

Remark. The results of this paper have been obtained independently by Goulden and Jackson [3]. We refer the reader to this very interesting paper for detailed applications and algorithms.

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References

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